

Kate's problem: The students are distributed around a circular table. The teacher distributes candies to all the students, so that each student has an even number of candies.

The transition: Each student passes half his store of candies to the right. After that is done, students with an odd number of candies eat one.

This game ends in finite time, with all students having the same number of candies.

Why? Let the number of students be N and the maximum stock of candies M , and let k be the number of students with M candies.

Theorem: If $k < N$, then either M decreases or k decreases.

Assuming $k < N$, there is one student with M candies such that the person to the left has M' candies with $M' < M$. Then after the move, that student has $M/2 + M'/2$ candies, if even, or $M/2 + M'/2 - 1$ otherwise; in either case, fewer candies. A student with A candies has at most $(A + M)/2$ candies after the move, and that is less than M , unless $A = M$. So nobody new has M candies. Thus, if $k = 1$, M decreases, otherwise k decreases.

So, there are at most $M(N - 1)$ moves that can be made before arriving at the situation where all students have the same number of candies. In this game, 0 is a possibility - but try to find an example!

A construction of the Icosahedron. Take three congruent golden rectangles. Call the line joining the midpoints of the shorter sides the "major axis". Place the rectangles in space so that they have the same center and the three major axes are mutually orthogonal. Now the vertices of the rectangles (12 in all) are the vertices of the icosahedron.

A Way of Looking at Fibonacci Numbers

Utah Teachers' Circle, September 24, 2007

A **recursive algorithm** is a procedure determined by two parts:

- A. An initial position;
- B. A rule for getting from one position to the next.

Example:

- A. Start at 1;
- B. Given a number, add 1.

It is plain to see that this gives us the counting numbers $1, 2, 3, \dots$ as far as we want to go. We can say that, for this algorithm the n th output is n .

Another example:

- A. Start at 1;
- B. At the n th step, add n .

Here, the n th number is $1 + 2 + 3 + \dots + n$. Can you show that this is $\frac{1}{2}n(n + 1)$?

Let a_n be the n th number. Draw a square $(n + 1) \times (n + 1)$ array of small squares. There are $n + 1$ squares on the diagonal, a_n square both above and below the diagonal. Thus

$$(n + 1)^2 = a_n + (n + 1) + a_n ,$$

hence the result. Or, proceed by induction which works since

$$\frac{1}{2}(n - 1)(n) + n = \frac{1}{2}n(n + 1) .$$

Another example:

- A. Start at 1;
- B. At the n th step, multiply by n .

Here, the n th number is $n!$, which of course is a symbol for what we have just done; the only way to calculate $n!$ is to multiply the result at hand by the next integer.

More examples:

(A) a_0, a_1 any two numbers ; $a_n = (a_{n-1} + a_{n-2})/2$;

Here the sequence is clearly bounded; in fact each term is between its two predecessors. This suggests convergence. Assuming convergence, can you find the limit?

Here's the trick: Find an r such that the sequence r^n satisfies the recursion relation:

$$r^2 = \frac{1}{2}(r + 1) ,$$

which has the roots

$$r = \frac{1 \pm \sqrt{8}}{4} = 1, -\frac{1}{2} .$$

So the sequences $\{1, 1, 1, \dots$ and $\{-.5, +.25, -.125, \dots, (-.5)^n, \dots$ both satisfy the recursion relation. So does any linear combination of these. So find coefficients α, β such that

$$a_0 = \alpha + \beta\left(\frac{1}{2}\right)^0 , \quad a_1 = \alpha + \beta\left(\frac{1}{2}\right)^1 .$$

We get this closed form:

$$a_n = \frac{a_0 + 2a_1}{3} + \frac{2}{3}(a_0 - a_1)\left(-\frac{1}{2}\right)^n .$$

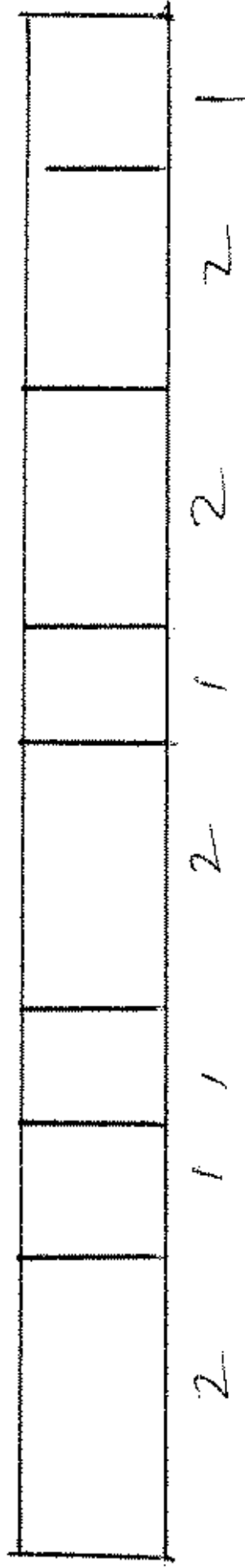
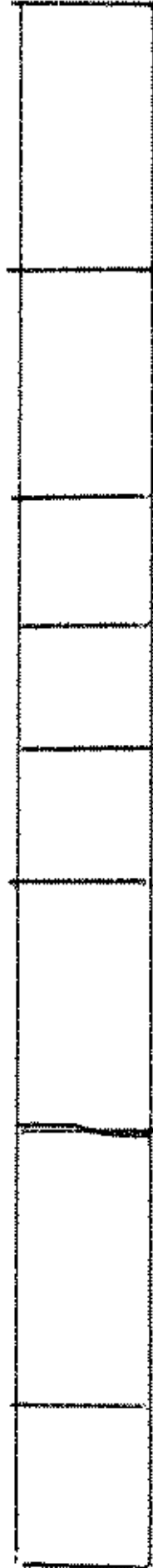
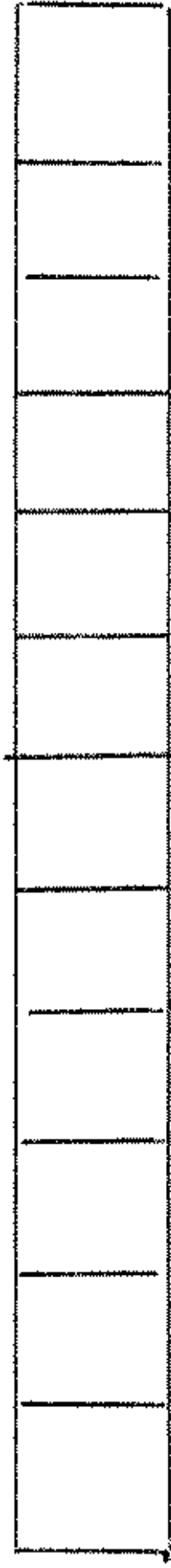
which clearly converges to $(a_0 + 2a_1)/3$.

$$(B) \quad a_0 \text{ any number ; } a_n = 1 + \frac{1}{a_{n-1}} .$$

We'll return to this sequence later. In the meantime, think about what its limit can be.

Now, let's start a new idea. Consider a line of n squares and piles of two types of tiles: 1-tiles are squares the size of the squares forming the line, and 2-tiles are two such squares with a common side.

Question: In how many ways can the line be covered by such tiles?



Let f_n be the number of ways to cover a line of length n . We can calculate the first few values of f_n .

$f_1 = 1$, obviously.

$f_2 = 2$: we can cover with two 1-tiles or 1 2-tiles.

$f_3 = 3$: there is one way using no 2-tiles, and two ways using 1 2-tile.

$f_4 = 5$: no 2-tiles: 1 way; one 2-tile: 3 ways; two 2-tiles: 1 way.

And so forth : $f_5 = 8, f_6 = 13, \dots$

Can you verify a recursion relation for this sequence?

Identity 1
$$f_1 = 1, \quad f_n = f_{n-1} + f_{n-2} .$$

Given a tiling, it ends in either a 1-tile or a 2-tile. We get f_{n-1} tilings that end in a 1-tile and f_{n-2} tilings that end in a 2-tile.

This is the **Fibonacci sequence**, and a wonderful way of describing it, leading to relatively easy combinatorial ways of deriving identities. For example, let's try to generalize the way we counted the first few values to get another formula for f_n .

Given n , in how many ways can we tile the n -line using i 2-tiles?

Hint: you'll need to know the combinatorial coefficient "n-choose-m" $\binom{n}{m}$.

Answer: Since each 2 tile covers two squares, the total number of 1-tiles used will be $n - 2i$, so the total number of tiles used is $n - i$. We can then conceive of the line as formed of $n - i$ spaces, of which i are to filled with 2-tiles and the remainder with 1-tiles. So the answer is the number of ways to select i spaces from $n - i$ spaces. This is

$$\binom{n-i}{i} = \frac{(n-i)!}{(n-2i)!i!}$$

and so we see that

Identity 2
$$f_n = \binom{n}{0} + \binom{n-1}{1} + \binom{n-2}{2} + \binom{n-3}{3} + \dots ,$$

where \dots means we keep adding terms so long as they make sense ($i \leq n - 1$).

Let's now look for other ways of counting the number of tilings, f_n .

For each k , suppose the last 2-tile covers squares $k + 1$ and $k + 2$. What identity does this lead to?

We cover the squares up to (and including) the k th square any way we please, but after the $k + 2$ th square, they are all 1-tiles. Thus there are f_k tilings with the last 2-tile in position $k + 1; k + 2$. We only go up to f_{n-2} , and finally there is the tiling with no 2-tiles. Thus

Identity 3
$$f_n = f_0 + f_1 + f_2 + \dots + f_{n-2} + 1 .$$

Now suppose the last 1-tile is in the $k + 1$ position. There are f_k such tilings, since every tile thereafter is a 2-tile. In particular, we must have $n - (k + 1)$ even, and so we better break this up into two cases.

n is odd. In this case, since $n - (k + 1)$ must be even, k can only be an even number. Thus

Identity 4
$$f_{2n+1} = f_0 + f_2 + f_4 + \cdots + f_{2n} .$$

n is even. Then we can only have k odd, and besides there is one tiling with no 1-tiles, so

Identity 5
$$f_{2n} = f_1 + f_3 + f_5 + \cdots + f_{2n-1} + 1 .$$

Now, put an n -line directly after an m -line and count the $m + n$ tilings where there is a break at this junction, and those where a 2-tile bridging the two lines. What identity does this produce?

There are $f_m f_n$ tiles with a break as described, and $f_{m-1} f_{n-1}$ tiles with no break there. Thus

Identity 6
$$f_{m+n} = f_m f_n + f_{m-1} f_{n-1} .$$

In particular

Identity 7
$$f_n^2 = f_n^2 + f_{n-1}^2 .$$

Let's start again with f_{m+n} , and iterate the recursion relation n times. What do you get?

$$\begin{aligned} f_{m+n} &= f_{m+n-1} + f_{m+n-2} \\ f_{m+n} &= f_{m+n-2} + f_{m+n-3} + f_{m+n-3} + f_{m+n-4} \quad \text{or} \\ f_{m+n} &= f_{m+n-2} + 2f_{m+n-3} + f_{m+n-4} \\ f_{m+n} &= (f_{m+n-3} + f_{m+n-4}) + 2(f_{m+n-4} + (f_{m+n-5})) + (f_{m+n-5} + f_{m+n-6}) \quad \text{or} \\ f_{m+n} &= f_{m+n-3} + 3f_{m+n-4} + 3f_{m+n-5} + f_{m+n-6} . \end{aligned}$$

If we do this again we start with f_{m+n-4} , and going down by one, the coefficients will be 1,4,6,4,1. In other words, the recursion is the same as that for Pascal's triangle, and in n steps we end up with

Identity 8
$$f_{m+n} = f_m + {}_1 C_n f_{m-1} + {}_2 C_n f_{m-2} + \cdots + {}_{n-1} C_n f_1 + f_0 .$$

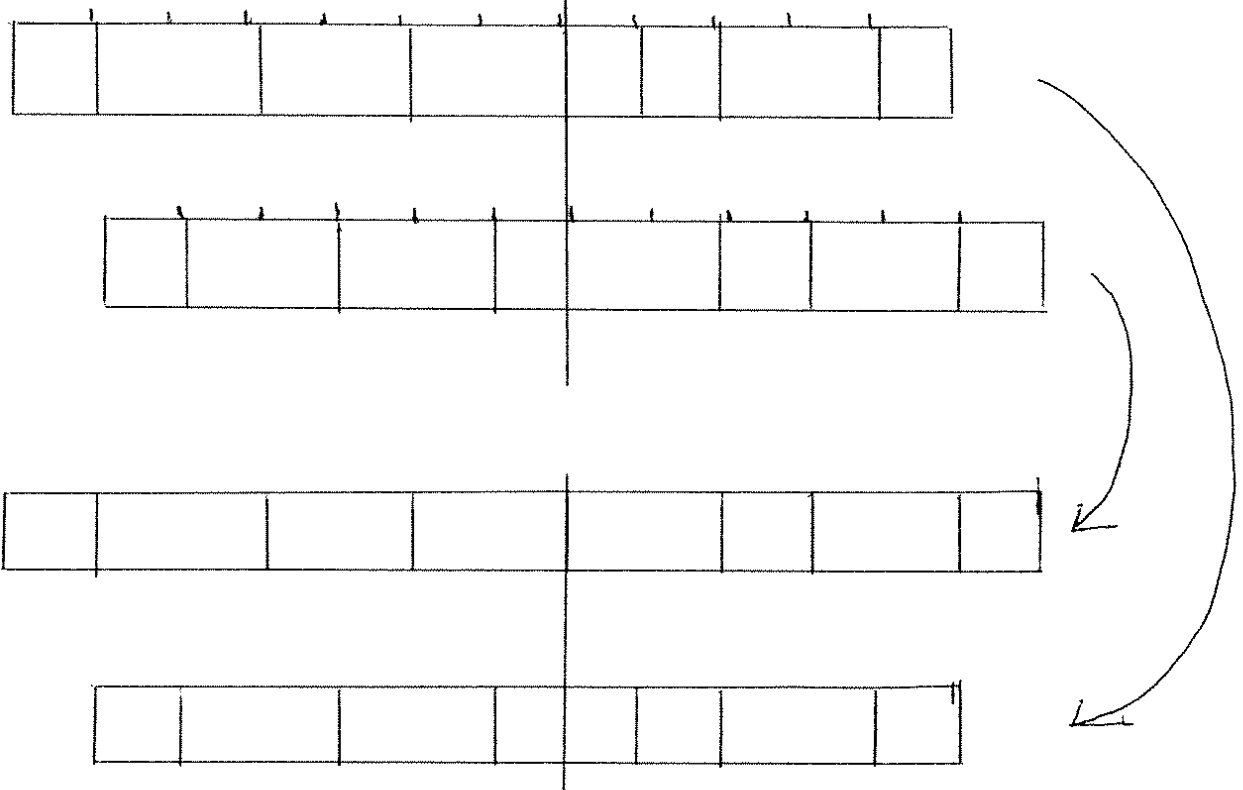
We return to the sequence defined recursively by

$$a_1 = 1; \quad a_n = 1 + \frac{1}{a_{n-1}} ,$$

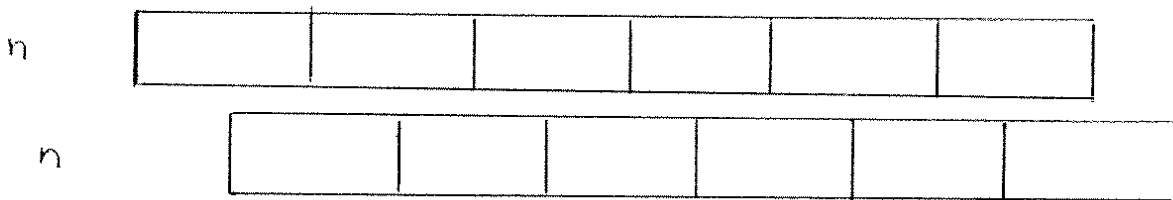
and calculate a few values:

$$1, 2, \frac{3}{2}, \frac{5}{3}, \frac{8}{5}, \frac{13}{8}, \frac{21}{13}, \frac{34}{21}, \frac{55}{34}, \frac{89}{55} .$$

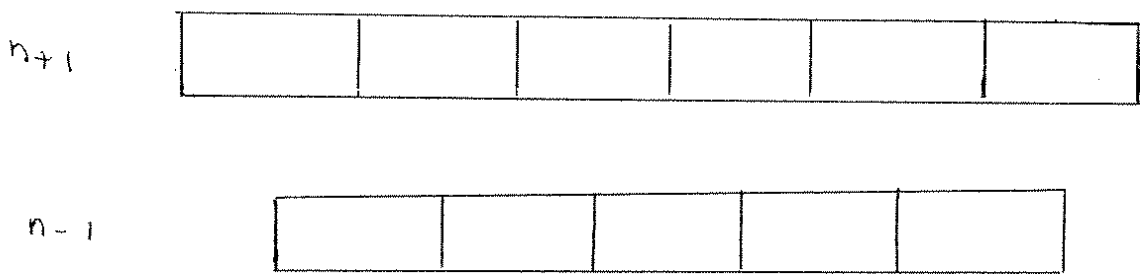
First
Fault



n even: only faultless



n odd: only faultless



We seem to be getting the quotients of consecutive members of the Fibonacci sequence. Can you show this, i.e., that

$$a_n = \frac{f_n}{f_{n-1}} .$$

Show that the quotients of consecutive fibonaccis satisfies the recursion relation and the initial condition:

$$\frac{f_n}{f_{n-1}} = \frac{f_{n-1} + f_{n-2}}{f_{n-1}} = 1 + \frac{f_{n-2}}{f_{n-1}} = 1 + \frac{1}{f_{n-1}/f_{n-2}} .$$

Now let's write the sequence in decimal form:

1, 2, 1.5, 1.66667, 1.6, 1.625, 1.61538, 1.61905, 1.61764, 1.61818, 1.61797, 1.61806, 1.61803 .

It looks like the sequence is settling down, somewhere around 1.6180. This suggests two things: that we show that the sequence actually converges, and find the limit. Assuming the first, can you find the limit?

Well, we can let $n \rightarrow \infty$ in the recursion equation, getting, where a represents the limit,

$$a = 1 + \frac{1}{a} .$$

This leads to the quadratic equation $a^2 - a - 1 = 0$ which has the roots

$$r = \frac{1 \pm \sqrt{5}}{2} .$$

Since $a \geq 1$ (why?), we must have

$$r = \frac{1 + \sqrt{5}}{2} ,$$

which is 1.618033989

So, the question remains: why does the sequence converge? That comes out of another identity that can be proven using the idea of tilings of lines of squares. Let's take two lines of length n . Then the number of different ways of tiling the two simultaneously is clearly f_n^2 . But now, let's count that another way. Place the second line just below the first line, but offset by just one square. Given a tiling of the two together, we'll say that it has a fault at square k (counting along the top line) if a new tile starts at square k on both lines.

Now, given a tiling, suppose it has faults and the first one is at square k . Exchange the tails of the two lines, getting a tiling of two lines, the top one of length $n + 1$, and the bottom one of length $n - 1$. This gives a correspondence of of all tilings-with-faults of two lines of length n with all tilings-with-faults of two lines of lengths $n - 1$ and $n + 1$. Finally, are there any tilings without faults in either case? We have to consider the even and odd cases separately.

If n is even, it is easy to see that the only no-fault tiling of two n lines (so arranged is that with all 2-tiles. Furthermore, there is no fault-free tiling of the two lines of lengths $n - 1$ and $n + 1$. Thus

$$n \text{ even : } f_n^2 = f_{n+1}f_{n-1} + 1 .$$

In case n is odd, each line must have at least 1 1-tile, from which we can conclude that all tilings have a fault. However, the two lines of lengths $n-1$ and $n+1$ have a no-fault tiling with all 2-tiles. Thus

$$n \text{ odd : } f_n^2 = f_{n+1}f_{n-1} - 1 ,$$

or, put together:

$$\text{Identity 9} \quad \text{forall } n : f_n^2 = f_{n+1}f_{n-1} + (-1)^n .$$

This then gives us

$$\text{Identity 10} \quad \frac{f_n}{f_{n-1}} = \frac{f_{n+1}}{f_n} + \frac{(-1)^n}{f_n f_{n-1}} .$$

One last point. The equation for this sequence is the same as the one we'd get for the fibonacci recursion, trying the sequence $a_n = r^n$. We conclude, in the same way as we did for the sequence of successive averages, that the fibonacci sequence has to be a linear combination of the roots $r = \frac{1 \pm \sqrt{5}}{2}$. If we fit such a linear combination to the initial values, we obtain

$$\text{Identity 11} \quad f_n = \frac{1}{\sqrt{5}} \left[\left(\frac{1 + \sqrt{5}}{2} \right)^n + \left(\frac{1 - \sqrt{5}}{2} \right)^n \right] .$$

Art Benjamin's magic trick: ask the participants to pick two integers and put them in a column. In the third row, write the sum of the two preceding rows, and continue, up to the tenth row. Now sum all the rows.

The sum is 11 times the 7th row.

Example

3
7
10
17
27
44
71*
115
186
301
sum: 781